# Solution procedures for the generalized discrete (r|p)centroid problem

D.R. Santos-Peñate<sup>1</sup>

C. Campos-Rodríguez<sup>2</sup> and J.A. Moreno-Pérez<sup>2</sup>

1. Dpto de Métodos Cuantitativos en Economía y Gestión/InstitutoUniversitario de Turismo y Desarrollo Económico Sostenible (Tides), Universidad de Las Palmas de Gran Canaria, Spain drsantos@dmc.ulpgc.es

> 2. Instituto Universitario de Desarrollo Regional, Universidad de La Laguna, Spain {campos, jamoreno}@ull.es

**Keywords** : competitive location, sequential location model, medianoid, centroid, Stackelberg, linear programming, metaheuristic, combinatorial optimization.

### **1** Introduction

In this work we consider a discrete competitive location model which is called (r/p)-centroid problem, leader-follower problem or Stackelberg problem in locations. The model represents a situation where two players, the leader and the follower make decisions sequentially in order to reach certain objectives. Leader and follower want to determine the locations for p and r facilities respectively. The objective of the follower, who makes the decision once the leader has selected its locations, is to maximize the demand captured by its facilities. The objective of the leader is to minimize the maximum demand that the follower could capture, as demand is assumed to be essential, this objective is equivalent to maximize the demand captured by its facilities. We study a generalized model in which the customer's choice rule is defined using a non-increasing capture function, the demand captured by the players is given by the value of this function for the difference between the distance from the demand point to the follower and the distance from the demand point to the leader. Given the set of p locations for the leader,  $X_p$ , the solution for the follower is an  $(r/X_p)$ -medianod. The solution for the leader is an (r/p)-centroid. We present a linear programming formulation for the generalized (r/p)-centroid problem and analyse some exact and heuristic solution procedures. We show some examples and results considering different capture functions.

## 2 Statement of the model and linear formulations

Let V be a set of points and C, L be subsets of V with cardinality |C|=m and |L|=n. C is the set of demand points, or clients, and L is the set of potential locations for facilities. Let  $d_{ij} = d(c_i, l_j)$  be the distance between demand point  $c_i$  and potential location  $l_j$ . For any  $c_i$  in C and any subset X of L,  $d_i(X)$  denotes the distance between  $c_i$  and X. Every point  $c_i$  in C has a weight  $w_i$  which represents the demand at point  $c_i$ . We consider a market of essential goods, which means that the demand is totally satisfied, that is, the sum of demands served by the firms operating in the market is equal to the total existing demand. The leader wants to open p facilities and the follower plans to enter the market with r facilities, both players want to determine the locations that allows them to reach certain objectives. The objective of the leader is to minimize the maximum demand that the follower could capture. Since demand is essential, this objective is equivalent to maximize the demand captured by the leader. The objective of the follower is to maximize the demand captured by its facilities.

Given the sets X and Y of locations for the leader's facilities and follower's facilities respectively, the demand captured by the follower is  $W = \sum_{i=1}^{n} w_i f_i(d_i)$  where  $d_i = d_i(Y) - d_i(X)$  and  $f_i$  is a non-increasing real function with  $0 \le f_i(d) \le 1$ , for any demand point  $c_i$  and  $d \ge 0$ . If  $\{Y_i\}_{i \in I}$  is the set of feasible solutions for the follower, the problem of the leader or (r/p)-centroid problem can be formulated as follows:

$$\begin{array}{l} \min W \\ & \prod_{j=1}^{n} x_{j} = p \\ & \sum_{i=1}^{m} \sum_{j=1}^{n} h_{ij}^{k} u_{ij} \leq W \quad k \in \left[1..\binom{n}{r}\right] \\ & \sum_{j=1}^{m} u_{ij} = 1 \qquad i \in [1..m] \\ & u_{ij} \leq x_{j} \qquad i \in [1..m], \ j \in [1..n] \\ & u_{ij}, \ x_{j} \in \{0,1\}, \ i \in [1..m], \ j \in [1..n] \end{array}$$

where  $x_j=1$  if the leader opens a facility at point  $l_j$  and  $x_j=0$  otherwise,  $u_{ij}=1$  if clients at  $c_i$  use a leader facility at point  $l_j$  and  $u_{ij}=0$  otherwise, and  $h_{ij}^k = w_i f_i (d_i (Y_k) - d_{ij})$ . Constraint  $u_{ij}=0,1$ , may be relaxed and replaced by  $u_{ij} \ge 0$ . The amount  $h_{ij}^k$  is the demand at  $c_j$  captured by the follower with the solution  $Y_k$  when the leader has a facility opened at  $l_j$  and  $l_j$  is the closest leader's facility to that demand point.

## 3 Solution procedures for the generalized (r/p)-centroid problem

Some exact and heuristic procedures proposed to solve the binary (r|p)-centroid problem (Alekseeva et al, Biesinger et al., Campos-Rodríguez et al., Roboredo and Pessoa) can be adapted to solve the generalized version. Considering different capture functions (binary, piecewise linear, piecewise concave and piecewise convex), we consider exact and heuristics solution approaches to solve the bi-level location problem. Some results are presented.

#### References

[1] E. Alekseeva, Y. Kochetov (2013). Matheuristics and exact methods for the discrete (r|p)-centroid problem. In: Talbi EG (ed) Metaheuristics for Bi-level Optimization. Studies in Computational Intelligence, vol. 482, 189-219. Springer Berlin Heidelberg.

[2] O. Berman, Z. Drezner, D. Krass (2010). Generalized coverage : New developments in covering location models. Computers & Operations Research 37, 1675-1687.

[3] O. Berman, D. Krass (2002). The generalized maximal covering location problem. Computers & Operations Research 29, 563-581.

[4] B. Biesinger, B. Hu, G. Raidl (2014). A hybrid genetic algorithm with solution archive for the discrete (r|p)-centroid problem. Technische Universitat Wien, Technical Report TR-186-1-14-03, 1-33.

[5] C.M. Campos-Rodríguez, D.R. Santos-Peñate, J.A. Moreno-Pérez (2010). An exact procedure and LP formulations for the leader-follower problem. TOP 18(1), 97-121.

[6] M.C. Roboredo, A.A. Pessoa (2013). A branch –and-cut algorithm for the discrete (r|p)-centroid problem. European Journal of Operational Research 224, 101-109.